

# BLOCK LU-FACTORIZATION OF HANKEL MATRICES

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# Outline

- 1 The origin
- 2 Block diagonalization of Hankel matrices
- 3 Hankel and Bezout matrices of two polynomials
- 4 Berlekamp–Massey
- 5 Bibliografy

The origin.

## To understand the Proof of Frobenius' Theorem

### Frobenius' Theorem

Let  $h$  be a regular Hankel matrix and  $h_{k_0} := 1$ .

Let  $h_{k_1}, \dots, h_{k_q}$  be the non-zero principal leading minors of  $h$   
( $1 \leq k_1 < \dots < k_q = n$ ).

Then the signature of  $h$  is given:

$$\text{Sig}(h) = \sum_{i=1}^q \begin{cases} 0 & \text{if } k_i - k_{i-1} \text{ is even} \\ (-1)^{\frac{r_i-1}{2}} \text{sign}(h_{k_i} h_{k_{i-1}}) & \text{if } k_i - k_{i-1} \text{ is odd} \end{cases}$$

F. Gantmacher. *Theory of Matrices, Vol. 1*, pag 343; Proof: 344–348.

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# Block diagonalization of Hankel matrices–Preliminaries

A Hankel matrix  $h$  is

$$h = \begin{pmatrix} \alpha_1 & \alpha_2 & \alpha_3 & \cdots & \alpha_n \\ \alpha_2 & \alpha_3 & & \ddots & \alpha_{n+1} \\ \alpha_3 & & \ddots & \ddots & \vdots \\ \vdots & \ddots & \ddots & & \vdots \\ \alpha_n & \alpha_{n+1} & \cdots & \cdots & \alpha_{2n-1} \end{pmatrix} = \mathbb{H}(\alpha_1, \dots, \alpha_{2n-1}),$$

# Block diagonalization of Hankel matrices–Preliminaries

Suppose that

$$h = H(\underbrace{0, \dots, 0}_{(r-1)}, \alpha_r, \dots, \alpha_{2n-1}) = \left( \begin{array}{ccc|ccc} 0 & \cdots & \alpha_r & \alpha_{r+1} & \cdots & \alpha_n \\ \vdots & \ddots & \vdots & \vdots & & \vdots \\ \alpha_r & \cdots & \alpha_{2r-1} & \vdots & & \vdots \\ \hline \vdots & & \vdots & \vdots & & \vdots \\ \vdots & & \vdots & \vdots & & \vdots \\ \alpha_n & \cdots & \alpha_{n+r-1} & \alpha_{n+r} & \cdots & \alpha_{2n-1} \end{array} \right)$$
$$= \left( \begin{array}{c|c} h_{11} & h_{12} \\ \hline h_{12}^t & h_{22} \end{array} \right),$$

# Block diagonalization of Hankel matrices–Preliminaries

$$\text{Then } h = \mathbb{H}(\underbrace{0, \dots, 0}_{(r-1)}, \alpha_r, \dots, \alpha_{2n-1}) = \begin{pmatrix} h_{11} & h_{12} \\ h_{12}^t & h_{22} \end{pmatrix},$$

and consider

$$T = \begin{pmatrix} \alpha_r & \cdots & & \alpha_{2n-1} \\ & \ddots & \ddots & \vdots \\ & & \ddots & \vdots \\ & & & \alpha_r \end{pmatrix} = \text{UpToep}(\alpha_r, \dots, \alpha_{2n-1});$$

$$t = \text{UpToep}(\alpha_r, \dots, \alpha_{n+r-1}).$$

# Block diagonalization of Hankel matrices

$$h = H(0, \dots, 0, \alpha_r, \dots, \alpha_{2n-1}), \quad T = \text{UpToep}(\alpha_r, \dots, \alpha_{2n-1}), \quad t = \text{UpToep}(\alpha_r, \dots, \alpha_{n+r-1})$$



# Block diagonalization of Hankel matrices

$h = H(0, \dots, 0, \alpha_r, \dots, \alpha_{2n-1})$ ,  $T = \text{UpToep}(\alpha_r, \dots, \alpha_{2n-1})$ ,  $t = \text{UpToep}(\alpha_r, \dots, \alpha_{n+r-1})$

**Theorem.** N. Ben-Atti et H. Lombardi

If  $T^{-1} = \text{UpToep}(\mu_1, \dots, \mu_{2n-r})$ ,  $t^{-1} = \text{UpToep}(\mu_1, \dots, \mu_n)$ , then

$$h' = (t^{-1})^t h t^{-1} = \begin{pmatrix} h'_{11} & 0 \\ 0 & h'_{22} \end{pmatrix},$$

where

$$h'_{11} = J h_{1,1}^{-1} J = \text{LowH}(\mu_1, \dots, \mu_r) \quad \text{and} \quad h'_{22} = -H(\mu_{r+2}, \dots, \mu_{2n-r}).$$

# Block diagonalization of Hankel matrices

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## Corollary

There is an Upper Triangular matrix  $A$  and a block diagonal matrix,  $D$ , such that

$$A^t h A = \text{Diag}(\text{lh}_{11}, \dots, \text{lh}_{tt})$$

and every  $\text{lh}_{ij}$  is a lower Hankel matrix.

# Block diagonalization–Euclidean Algorithm

Inversion of Toeplitz matrices: If

$$T = \text{UpToep}(\alpha_1, \dots, \alpha_{2n-1}), \alpha_1 \neq 0, T^{-1} = \text{UpToep}(\mu_1, \dots, \mu_{2n-1})$$

and we define

$$S(x) = \alpha_1 + \alpha_2x + \dots + \alpha_{2n-1}x^{2n-2}, Q(x) = \mu_1 + \mu_2x + \dots + \mu_{2n-1}x^{2n-1},$$

then  $S(x)Q(x) = 1 \pmod{x^{2n-1}}$

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This leads the Algorithm for Computing the blocks  $lh_{ij}$ : Let

- $R_0 = x^{2n-1}; R_1 = \alpha_1 x^{2n-2} + \dots + \alpha_{2n-1}$ ;
- $\{Q_1, \dots, Q_{t-1}\}$  : Quotients in the signed Euclidean algorithm for  $(R_0, R_1)$ ;

$$Q_i = c_{0,i} x^{d_1} + \dots + c_{d_1,i}$$

- $R_t$ : the last non zero remainder,

then

$$lh_{ij} = \text{LowH}(c_{0,i}, \dots, c_{d_{i-1},i}), i \leq t-1; lh_{tt} = \text{LowH}(\text{ from } R_t)$$

# Proof of Frobenius' Theorem

$$A^t h A = D = \text{Diag}(lh_{11}, lh_{22}, \dots, lh_{tt})$$

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We have

$$\text{Sig}(h) \stackrel{\text{Sylvester's law}}{=} \text{Sig}(D) = \text{Sig}(\text{Diag}(lh_{11}, \dots, lh_{qq})) = \sum_{i=1}^q \text{Sig}(lh_{ii}). \quad (1)$$

Since

$$\text{Sig}(lh_{ii}) = \begin{cases} 0 & \text{if } d_i \text{ is even,} \\ \text{sign}(c_{0,i}) = (-1)^{\frac{d_i-1}{2}} \text{sign}(\det(lh_{ii})) & \text{if } d_i \text{ is odd;} \\ = (-1)^{\frac{d_i-1}{2}} \text{sign}(\det(D_{k_i} D_{k_{i-1}})) & d_i = k_i - k_{i-1} \\ = (-1)^{\frac{d_i-1}{2}} \text{sign}(\det(h_{k_i} h_{k_{i-1}})) & \end{cases} \quad (2)$$

it follows the **Frobenius' Theorem** from (1) and (2)

$$\text{Sig}(h) = \sum_{i=1}^q \begin{cases} 0 & \text{if } k_i - k_{i-1} \text{ is even} \\ (-1)^{\frac{r_i-1}{2}} \text{sign}(h_{k_i} h_{k_{i-1}}) & \text{if } k_i - k_{i-1} \text{ is odd} \end{cases}$$

# Literature

In [Gragg-Lindquist] (1983) and [Heinig-Rost] (1984), they assert that a Hankel matrix is **LU-equivalent** to a block diagonal matrix.

## Theorem

There is an upper triangular matrix  $A_1$  with entries equal to one on the main diagonal such that

$$A_1^t h A_1 = D_1$$

where  $D_1$  is a block diagonal matrix and each block is a lower Hankel triangular matrix.

Such a block diagonalization can also be found for example in [Bini-Pan] and [Barel-Bultheel].

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# Hankel and Bezout matrices of two polynomials

Let  $u(x)$  and  $v(x)$  be two coprime polynomials,

$$u(x) = \sum_{i=0}^n u_i x^i; \quad v(x) = \sum_{i=0}^m v_i x^i, \quad n > m.$$

- The  $n \times n$  Hankel matrix associated to  $u(x)$  and  $v(x)$  is defined as

$$\mathbf{H}(u, v) = \mathbf{H}(h_1, h_2, \dots, h_{2n-1}),$$

where

$$R(x) = \frac{v(x)}{u(x)} = \sum_{i=1}^{\infty} h_i x^{-i}.$$

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where

$$R(x) = \frac{v(x)}{u(x)} = \sum_{i=1}^{\infty} h_i x^{-i}.$$

- The Bezout Matrix associated to  $u(x)$  and  $v(x)$  is the symmetric matrix

$$\mathbf{Bez}(u, v) = \begin{pmatrix} c_{0,0} & \dots & c_{0,n-1} \\ \vdots & & \vdots \\ c_{n-1,0} & \dots & c_{n-1,n-1} \end{pmatrix},$$

where the  $c_{i,j}$  are defined by the Cayley expression:

$$\frac{u(x)v(y) - u(y)v(x)}{x - y} = \sum_{i,j=0}^{n-1} c_{i,j} x^i y^j.$$

# Hankel and Bezout matrices of two polynomials

## PROPERTIES

- $\text{degree}(\gcd(u, v)) = n - \text{rk}(\text{Bez}(u, v)) = n - \text{rk}(H(u, v))$ .
- Both are symmetric.
- Both matrices provide subresultant polynomials.
- Both matrices represent the multiplication by  $v(x)$  modulo  $u(x)$ .
- $\text{Bez}(u, v) = \text{Bez}(u, 1)H(u, v)\text{Bez}(u, 1)$ , with

$$\text{Bez}(u, 1) = \text{UpperHk}(u_1, \dots, u_n).$$

$$\blacktriangleright J \text{Bez}(u, v) J = J \text{Bez}(u, 1) H(u, v) \text{Bez}(u, 1) J$$

$$\Downarrow H(u, v) = A^{-t} D A^{-1}$$

$$\star J \text{Bez}(u, v) J = [J \text{Bez}(u, 1) A^{-t}] D [A^{-1} \text{Bez}(u, 1) J]$$

# Block Diagonalization for $H(u, v)$

If we apply our algorithm, in the first step

$$(t^{-1})^t H(u, v) t^{-1} = \begin{pmatrix} \text{JBez}(q_1, 1) \text{J} & 0 \\ 0 & H(v, r_1) \end{pmatrix}$$

with

$$u(x) = v(x)q_1(x) - r_1(x), \quad t = \text{UpToep}(h_{n-m}, \dots, h_{2n-m}).$$

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In the last one,

$$A^t H(u, v) A = \begin{pmatrix} \text{JBez}(q_1, 1) \text{J} & & & \\ & \text{JBez}(q_2, 1) \text{J} & & \\ & & \ddots & \\ & & & H(r_{t-1}, r_t) \end{pmatrix}$$

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Moreover, the rows of  $\mathbf{A}^{-1} \text{Bez}(\mathbf{u}, \mathbf{1}) \text{J}$  are defined by the remainders.

# Block Diagonalization for Bez( $u, v$ )

In [Bini-Gemignani] (1995),

$$\begin{aligned} J\text{Bez}(u, v)J &= \begin{pmatrix} (J\text{Bez}(u, v)J)_{m_i} & X^t \\ X & W \end{pmatrix} \\ &= \begin{pmatrix} T & 0 \\ K & I \end{pmatrix} \begin{pmatrix} J\text{Bez}(u^{(i)}, v^{(i)})J & 0 \\ 0 & J\text{Bez}(r_{i-1}, r_i)J \end{pmatrix} \begin{pmatrix} T^t & K^t \\ 0 & I \end{pmatrix}, \end{aligned}$$

where

$$H(u, v)_{m_i} = H(u^{(i)}, v^{(i)}), \quad J\text{Bez}(r_{i-1}, r_i)J : \text{Schur complement of } (J\text{Bez}(u, v)J)_{m_i},$$

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If we sequentially apply the relation, then

$$A_B^t J\text{Bez}(u, v)J A_B = \begin{pmatrix} J\text{Bez}(q_1, 1)J & & & \\ & J\text{Bez}(q_2, 1)J & & \\ & & \ddots & \\ & & & J\text{Bez}(r_{t-1}, r_t)J \end{pmatrix}$$



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Moreover, the rows of  $A_B^{-1}$  are defined by the remainders.

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# Berlekamp–Massey

## Classical

**We have:** The first  $2n$  coefficients of a linearly recurrent sequence,  $[a_0, a_1, \dots, a_{2n-1}]$ . The minimal polynomial has degree bound  $n$ .

**We want to compute:** The minimal polynomial of the sequence  $P(x)$ .

INPUT:

$$R_0 := x^{2n}, R_1 := a_0 + a_1 x + \dots + a_{2n-1} x^{2n-1}; V_0 = 0; V_1 = 1;$$

GOAL:  $V$  such that  $\text{degree}(V R_1) < n$

ALGORITHM:

while  $n \leq \text{degree}(R_1)$  do

$(Q, R) :=$  quotient and remainder of  $R_0$  divided by  $R_1$  ;

$V := V_0 - Q V_1$  ;

$V_0 := V_1$  ;  $V_1 := V$  ;  $R_0 := R_1$  ;  $R_1 := R$  ;

end while;

$d := \max(\text{degree}(V_1), 1 + \text{degree}(R_1))$  ;  $\mathbf{P} := \mathbf{x}^d \mathbf{V}_1(\mathbf{1}/\mathbf{x})$  ;

# Berlekamp–Massey

Modified

**We have:** The first  $2n$  coefficients of a linearly recurrent sequence,  $[a_0, a_1, \dots, a_{2n-1}]$ . The minimal polynomial has degree bound  $n$ .

**We want to compute:** The minimal polynomial of the sequence  $P(x)$ .

INPUT:

$$R_0 := x^{2n}, R_1 := a_0x^{2n-1} + a_1x^{2n-2} + \dots + a_{2n-1}; V_0 = 0; V_1 = 1;$$

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$V := V_0 - Q V_1$  ;

$V_0 := V_1$  ;  $V_1 := V$  ;  $R_0 := R_1$  ;  $R_1 := R$  ;


**end while;**

**P :=  $V_1/\text{lc}(V_1)$  ;**


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**Merci Henri**

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